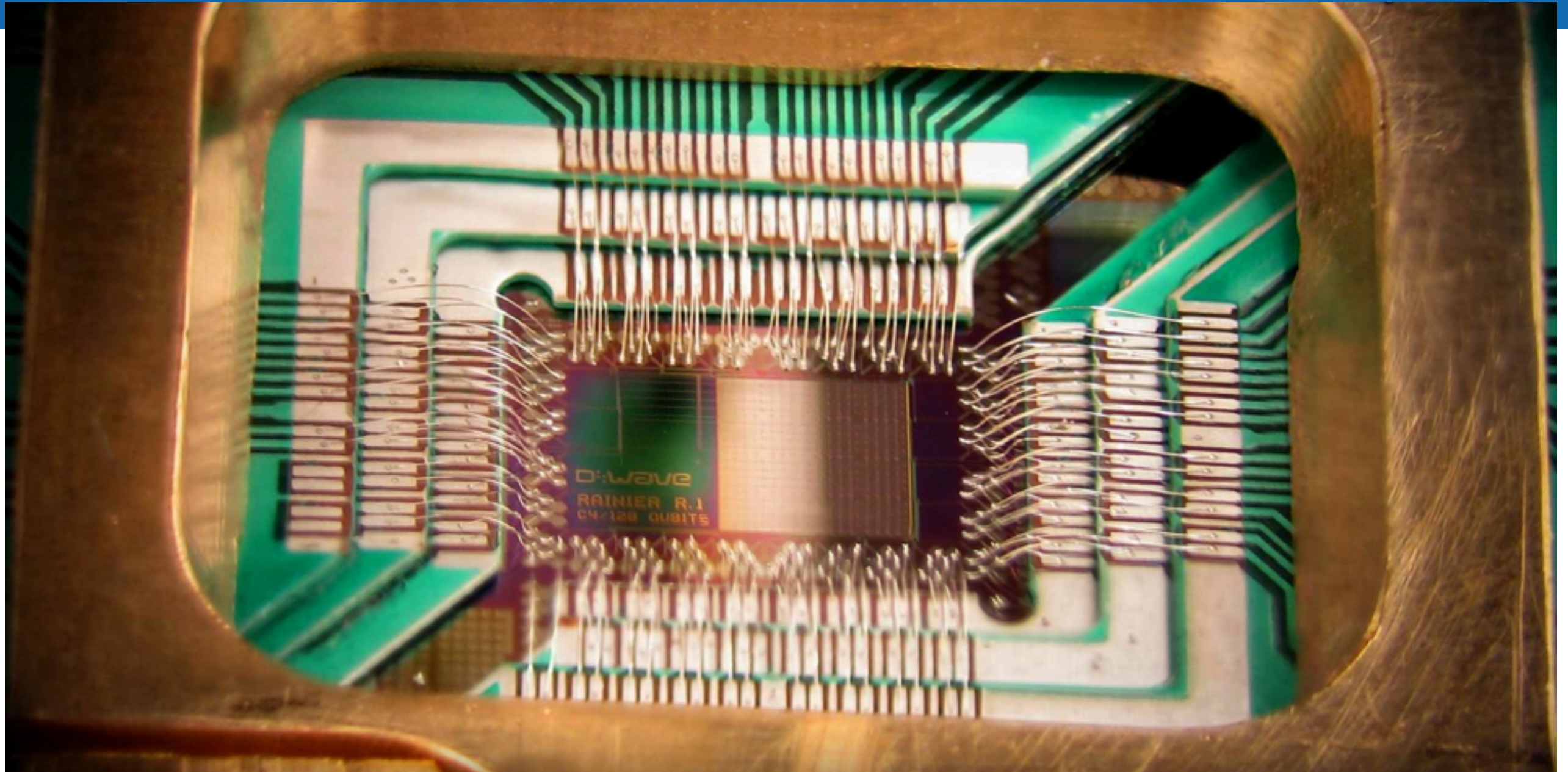




High Performance *Quantum* Computing

Matthias Troyer

Beyond exascale: quantum devices



Quantum random numbers
perfect randomness



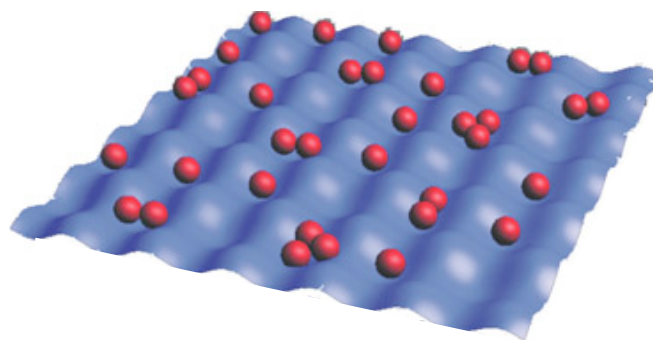
Quantum encryption
secure communication



Quantum computers?

- solve quantum models (R. Feynman)
- factor integers (P. Shor)
- break RSA encryption
- solve linear equations (A. Harrow *et al*)

...



Analog quantum simulators
solve quantum models



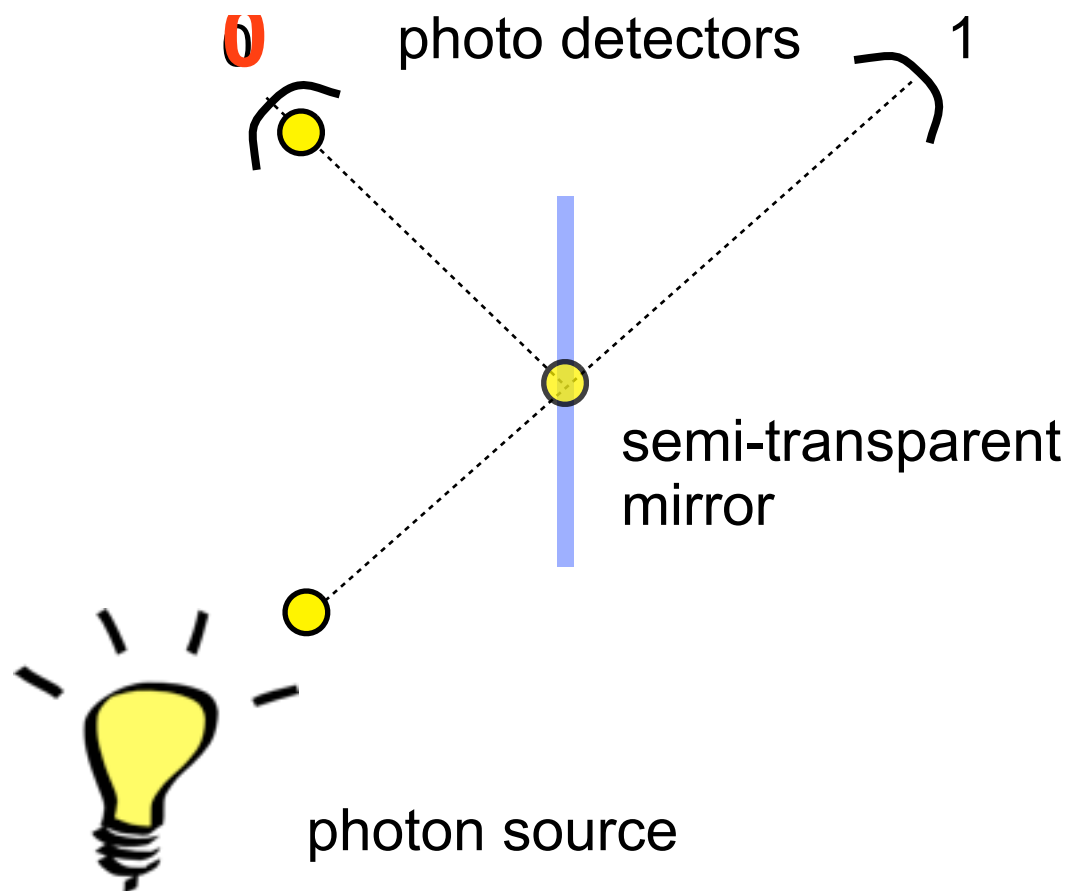
Quantum annealer
solve hard optimization problems?

Quantum random numbers

Quantum mechanics can give true and perfect random numbers



1. Photon source emits a photon
2. Photon hits semi-transparent mirror
3. Photon follows both paths
4. The photo detectors see the photon only in one place: **random selection**
5. Record one random bit



Quantum cryptography

Exchange a secret key to encrypt a message using entangled photons



No eavesdropper can listen without being detected

except for the quantum hacker Vadim Makarov (Waterloo)



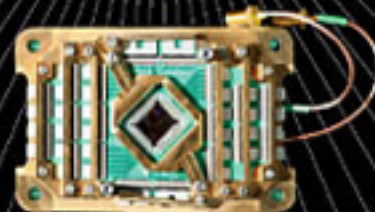
French Advances / My Doctor Fired Me / Love App-tually

TIME

IT PROMISES TO SOLVE SOME OF HUMANITY'S
MOST COMPLEX PROBLEMS. IT'S BACKED
BY JEFF BEZOS, NASA AND THE CIA.
EACH ONE COSTS \$10,000,000 AND OPERATES
AT 459° BELOW ZERO. AND NOBODY KNOWS
HOW IT ACTUALLY WORKS

THE INFINITY MACHINE

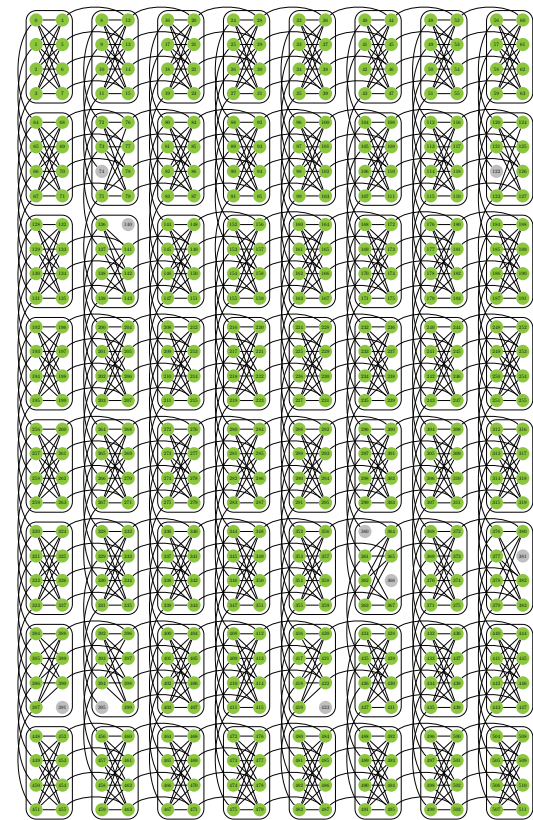
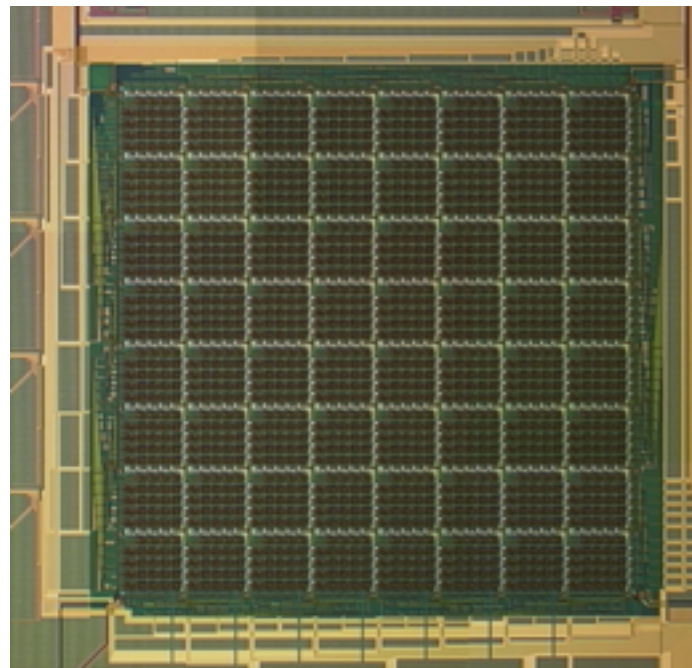
BY LEV GROSSMAN



D-Wave Two quantum annealer

Solves NP-complete spin glass problem with up to 512 variables

$$H = \sum_{ij} J_{ij} s_i s_j + \sum_i h_i s_i + \text{const.} \quad \text{with} \quad s_i = \pm 1$$

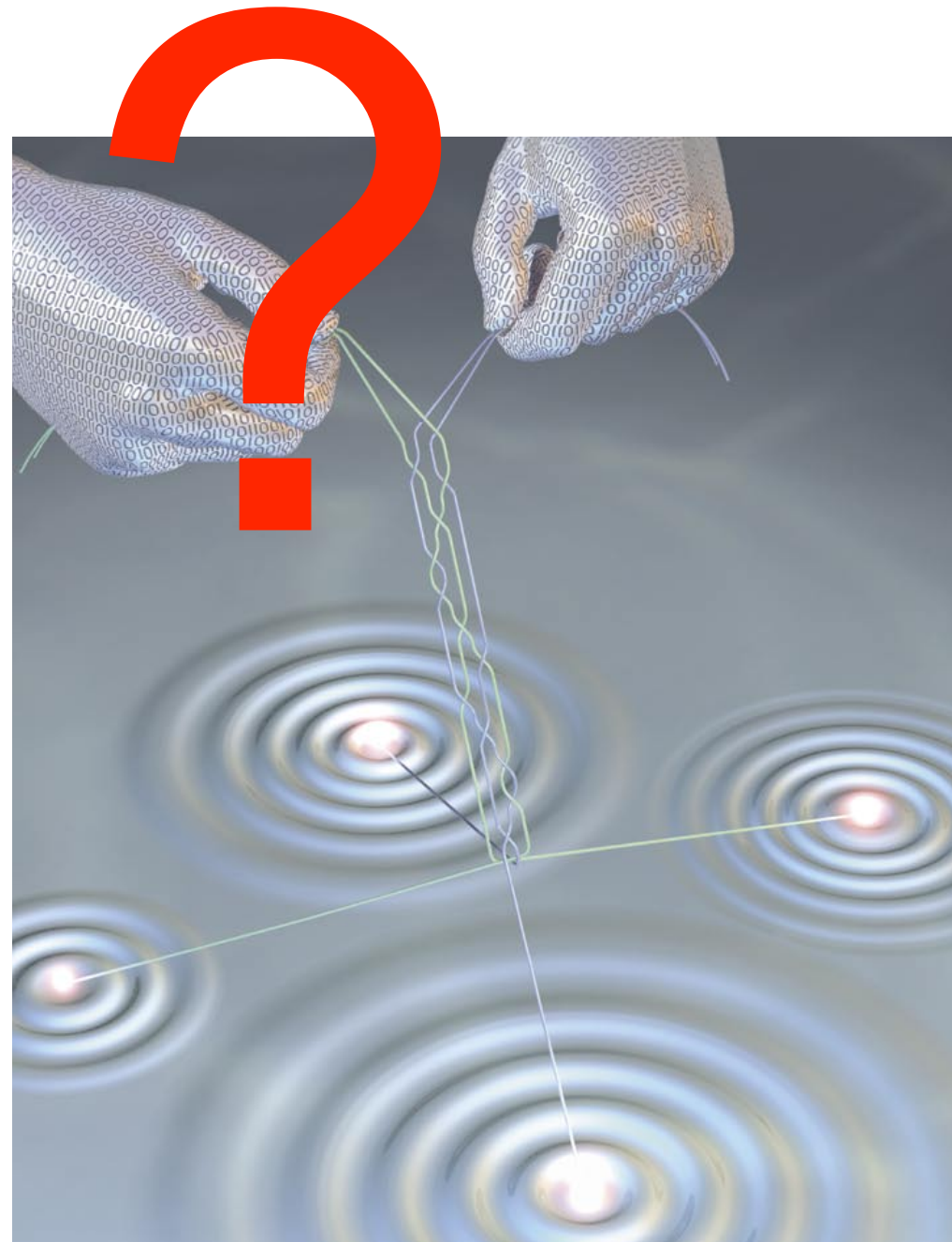


Can be built with imperfect qubits

Suffers from calibration problems like other analog devices

Unknown if this technology can ever scale better than classical devices

What about quantum computers ?



Microsoft
Windows Q

Classical versus quantum bits

- Classical bits can only be either 0 or 1
- Quantum bits can have both values at once, in arbitrary superpositions

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle \quad |\alpha|^2 + |\beta|^2 = 1$$

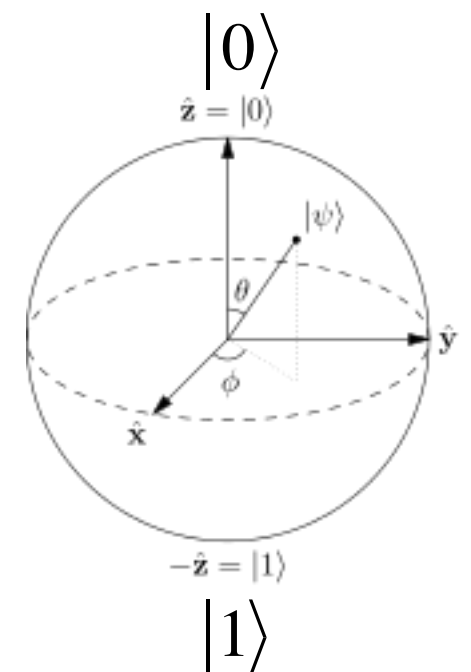
- need two complex or three real numbers to describe the state

$$|\psi\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \quad \alpha = \cos\left(\frac{\theta}{2}\right) \quad \beta = e^{i\phi} \sin\left(\frac{\theta}{2}\right)$$

- but when measuring (looking) I only ever get 1 bit

0 with probability $|\alpha|^2$

1 with probability $|\beta|^2$



Information content of a quantum register

- The state of an N qubit register
 - needs 2^N complex numbers to be represented classically
 - but when measured only gives N bit of information
- Advantages and limitations of quantum computers
 - Exponential intrinsic parallelism: operate on 2^N inputs at once
 - But very limited readout of only N bits
- No-cloning theorem: a quantum register cannot be copied

Try to build a
cloner for 0 and 1

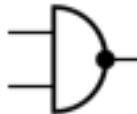
$$C|0\rangle \rightarrow |0\rangle|0\rangle$$

$$C|1\rangle \rightarrow |1\rangle|1\rangle$$

Linearity of quantum mechanics means it
will not clone an arbitrary state

$$\begin{aligned} C(\alpha|0\rangle + \beta|1\rangle) &\rightarrow \alpha|0\rangle|0\rangle + \beta|1\rangle|1\rangle \\ &\neq (\alpha|0\rangle + \beta|1\rangle)(\alpha|0\rangle + \beta|1\rangle) \end{aligned}$$

Quantum gates

- In classical computing the NAND gate is universal 
- In quantum computing we need three gates

- Hadamard gate (flips x and z)

$$\text{---} \boxed{H} \text{---} \quad |0\rangle \rightarrow \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \quad |1\rangle \rightarrow \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

- T gate (applies phase to 1 state)

$$\text{---} \boxed{T} \text{---} \quad \alpha|0\rangle + \beta|1\rangle \rightarrow \alpha|0\rangle + e^{i\pi/4}\beta|1\rangle$$

- CNOT gate (conditionally flips)

$$\begin{array}{c} \text{---} \bullet \text{---} \\ | \\ \text{---} \oplus \text{---} \end{array} \quad |x\rangle|y\rangle \rightarrow |x\rangle|x \oplus y\rangle$$

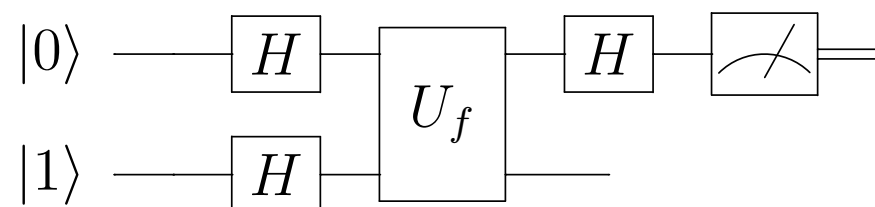
- This choice is not unique and having more gates can make a device more efficient

The Deutsch algorithm: simplest quantum speedup

- Check whether a binary function is constant or not

$$f : \{0,1\} \rightarrow \{0,1\}$$

- Classically two function calls are needed: $f(0) = f(1)$?
- Quantum mechanically only one function call by applying the function to both arguments at once



$$U_f |x\rangle |y\rangle \rightarrow |x\rangle |f(x) \oplus y\rangle$$

$$U_f (\alpha |0\rangle + \beta |1\rangle) |0\rangle \rightarrow \alpha |0\rangle |f(0)\rangle + \beta |1\rangle |f(1)\rangle$$

- Smart manipulation needed to read out the answer as one bit

$$\frac{1}{2} \left[\left(1 + (-1)^{f(0) \oplus f(1)} \right) |0\rangle + \left(1 - (-1)^{f(0) \oplus f(1)} \right) |1\rangle \right]$$

What would we use it for?

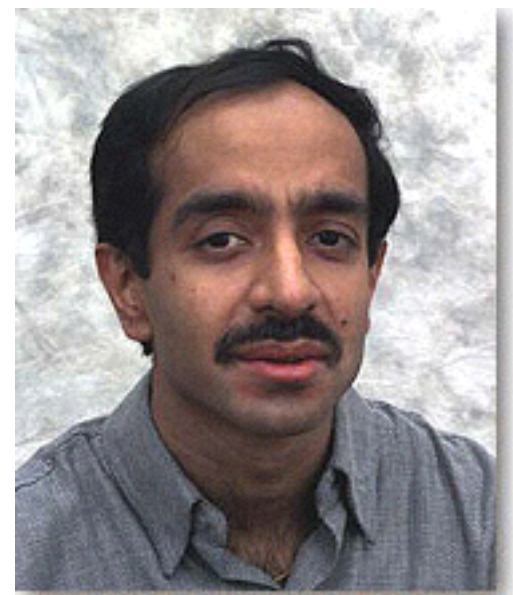
Surprisingly, this question has not been seriously asked until 2012

It needs to be an important problem (killer-app) that we cannot solve on an exascale machine

Compare to state-of-the-art on beyond-exascale hardware and not an unoptimized code on a single CPU core

Grover search

- Search an unsorted database of N entries in $\sqrt{N}$ time
- Rare case of provable quantum speedup given an oracle
- However, the oracle needs to be implemented!
 - N -entry database or arbitrary stored data needs at least $O(N)$ hardware
 - Can perform the same search classically in $\log(N)$ time with special purpose hardware
- Grover search is only useful if the database need not be stored but can be calculated on the fly!



Shor's algorithm for factoring

- Factoring is a hard problem classically: $O(\exp(N^{1/3}))$ for N bits

```

53693968364269119460795054153326005186041818389302311662023173188470613584169777981247775554355964649
04452615804209177029240538156141035272554197625377862483029051809615050127043414927261020411423649694
63096709107717143027979502211512024167962284944780565098736835024782968305430921627667450973510563924
  02989775917832050621619158848593319454766098482875128834780988979751083723214381986678381350567167
                                =
43636376259314981677010612529720589301303706515881099466219525234349036065726516132873421237667900245
9135372537443549282380180405548453067960658656053548608342707327969894210413710440109013191728001673
                                *
12304864190643502624350075219901117888161765815866834760391595323095097926967071762530052007668467350
6058795416957989730803763009700969113102979143329462235916722607486848670728527914505738619291595079

```

- Shor's algorithm is polynomial time on a quantum computer
 - $O(N^3)$ using minimal number of $2N+3$ qubits
 - $O(N^2)$ using $O(N)$ qubits
 - $O(N)$ using $O(N^2)$ qubits
- Shor's algorithm suddenly made quantum computing interesting

Shor's algorithm and encryption

- Shor's algorithm can be used to crack RSA encryption
 - assuming 10 ns gate time and minimal number of $2N+3$ qubits
 - much faster (seconds) when using more qubits

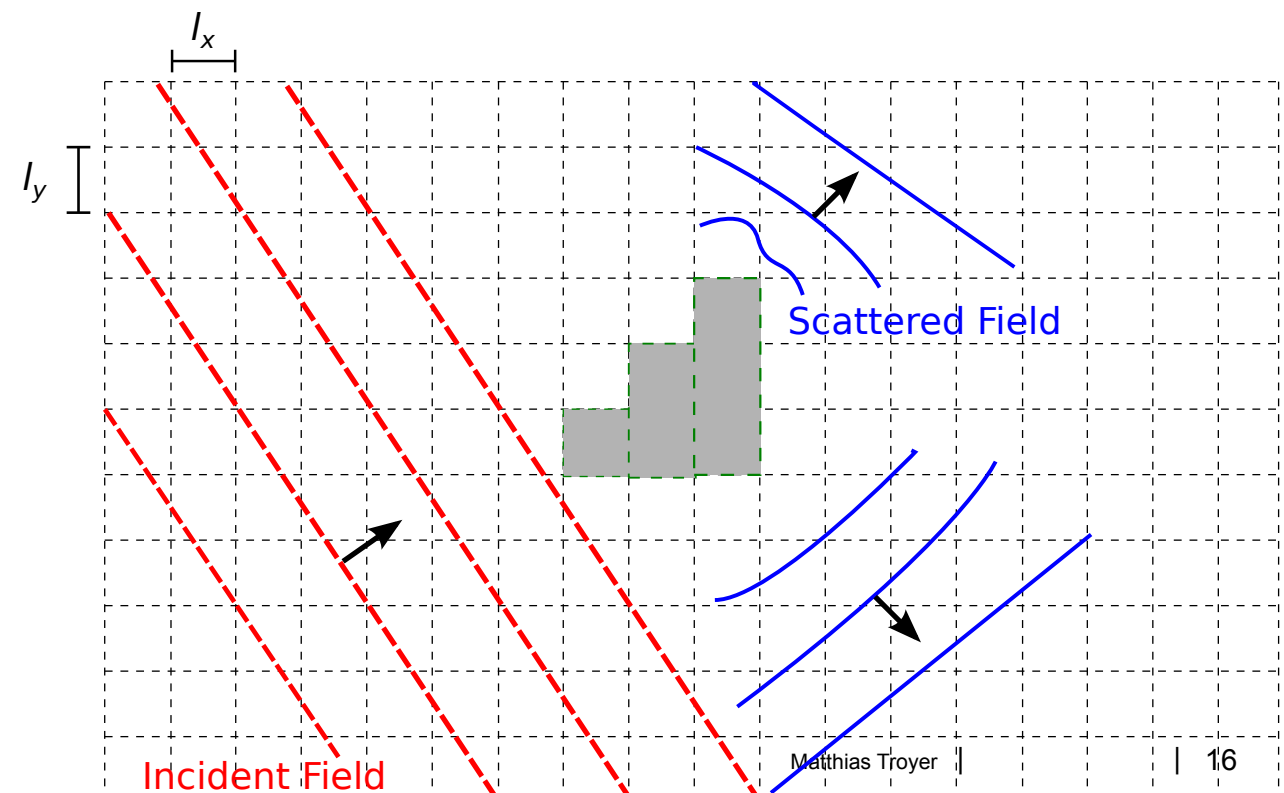
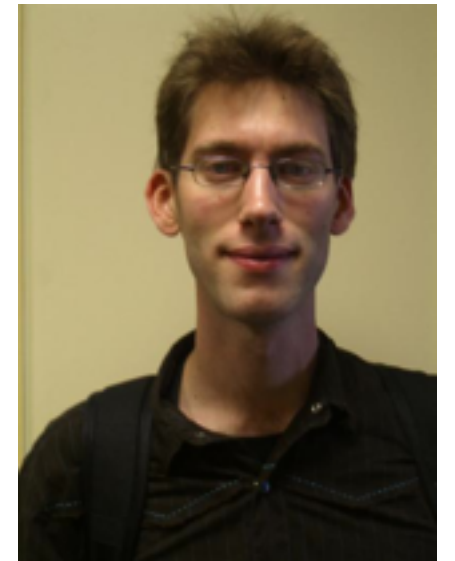
RSA	cracked in	CPU years	Shor
453 bits	1999	10	1 hour
768 bits	2009	2000	5 hours
1024 bits		1000000	10 hours



- But use of quantum computers to crack RSA is limited since we can anytime switch to post-quantum encryption
 - quantum cryptography
 - lattice based cryptography

Solving linear systems of equations (Harrow *et al*)

- Solve linear systems in $\log(N)$ time if
 - matrix can be computed efficiently and need not be stored
 - only $\log(N)$ bits of the answer are needed
 - problem is well conditioned
- First application: electromagnetic wave scattering
 - Clader, Jacobs, Sprouse (2013)
 - crossover compared to classical supercomputers is beyond 1000 years wallclock time



Applications running at scale on Jaguar @ ORNL

Domain area	Code name	Institution	# of cores	Performance	Notes
Materials	DCA++	ORNL	213,120	1.9 PF	2008 Gordon Bell Prize Winner
Materials	WL-LSMS	ORNL/ETH	223,232	1.8 PF	2009 Gordon Bell Prize Winner
Chemistry	NWChem	PNNL/ORNL	224,196	1.4 PF	2008 Gordon Bell Prize Finalist
Materials	DRC	ETH/UTK	186,624	1.3 PF	2010 Gordon Bell Prize Hon. Mention
Nanoscience	OMEN	Duke	222,720	> 1 PF	2010 Gordon Bell Prize Finalist
Biomedical	MoBo	GaTech	196,608	780 TF	2010 Gordon Bell Prize Winner
Chemistry	MADNESS	UT/ORNL	140,000	550 TF	
Materials	LS3DF	LBL	147,456	442 TF	2008 Gordon Bell Prize Winner
Seismology	SPECFEM3D	USA (multiple)	149,784	165 TF	2008 Gordon Bell Prize Finalist

Solving quantum chemistry on a quantum computer

- A killer-app for quantum computing is solving quantum problems
 - Design a room-temperature superconductor
 - Develop a catalyst for carbon sequestration
 - Develop better catalysts for nitrogen fixation (fertilizer)
- These problem need better accuracy than we get by using approximate classical algorithms
 - exponentially hard classically
 - polynomial complexity on quantum hardware



Quantum algorithms for quantum chemistry Microsoft Research

- Assume the fastest imaginable quantum computer and consider smallest problem we cannot solve classically: $N=50$ orbitals, $O(100)$ qubits

	Scaling	Run time for $N=50$
State of 12/2013	N^{10}	100 000 years
1/2014 pipelining and optimisations	N^9	3 years
6/2014 faster convergence	$N^{5.5}$	1 hour
parallel with with	$N^{1.6}$	seconds?

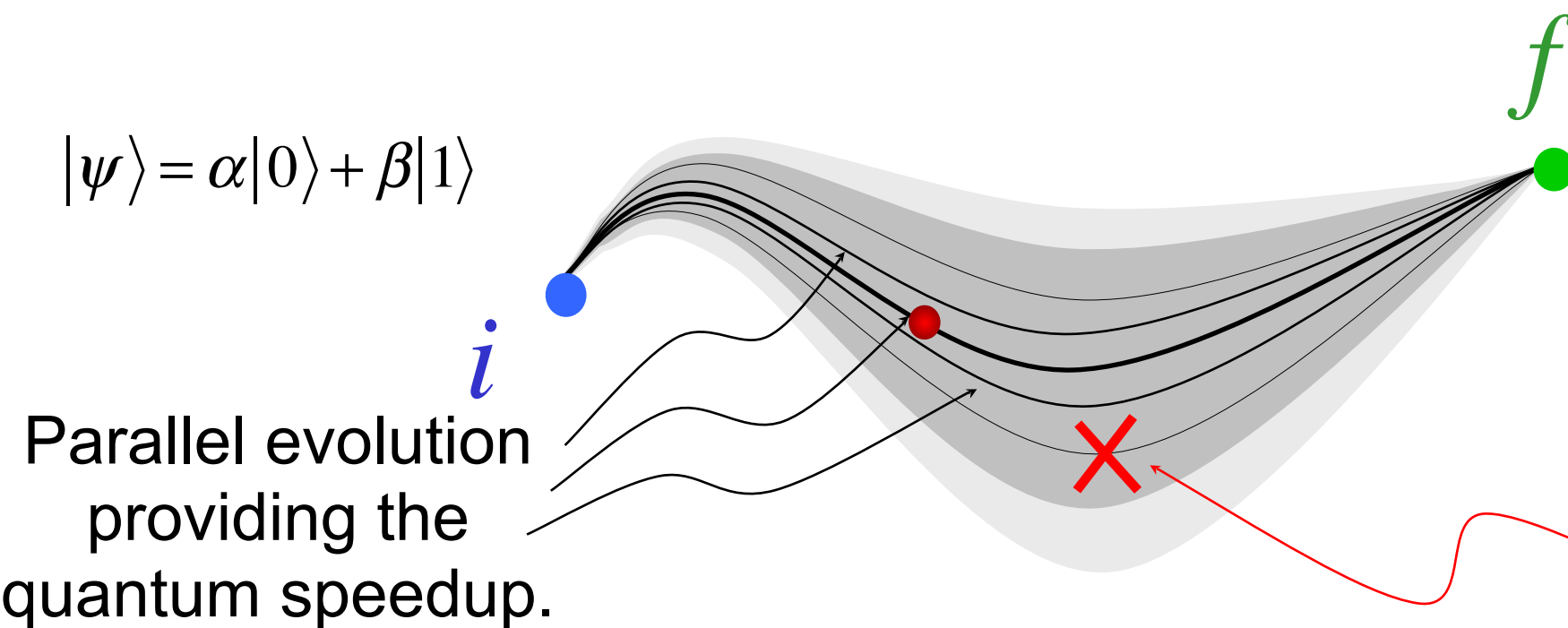
- Quantum information theorists declare victory proving the existence of polynomial time algorithms
- We need HPQC specialists to develop better algorithms!

State of the art of hardware

many proposals but few are scalable ...

Physical realizations and decoherence

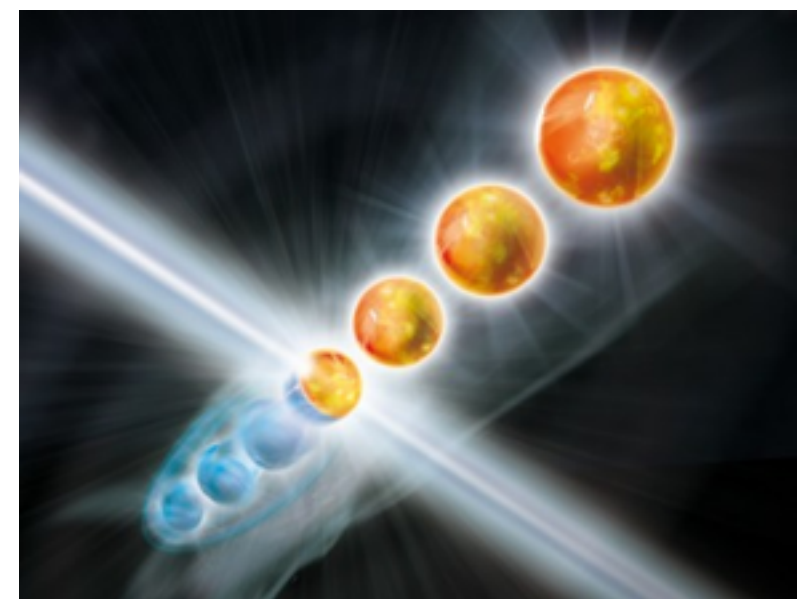
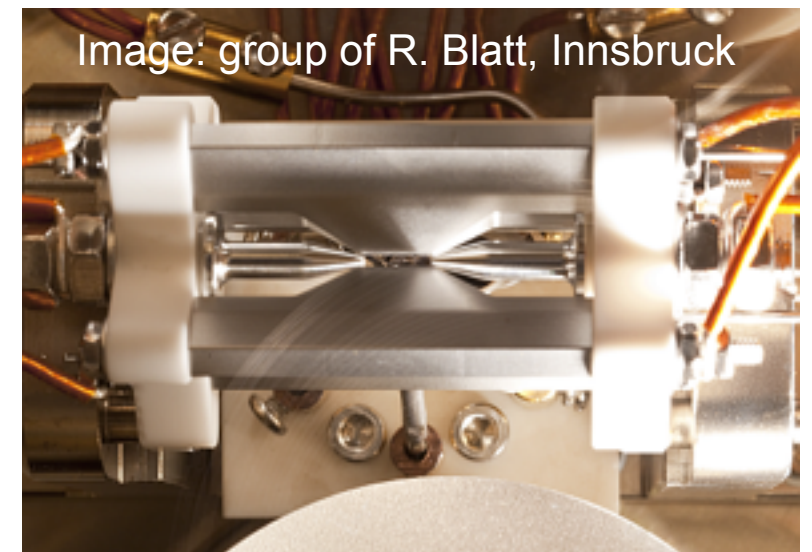
- Coupling to the environment easily destroys the quantum superposition
- Qubits need to be designed to be well isolated from the environment



Perturbations
from the environment
destroy the parallel evolution of the computation

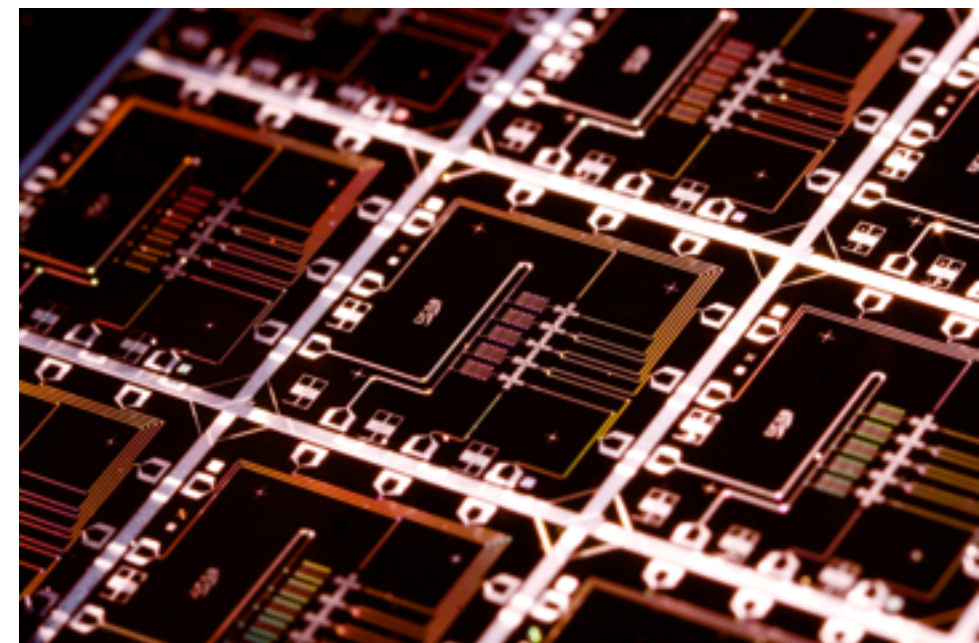
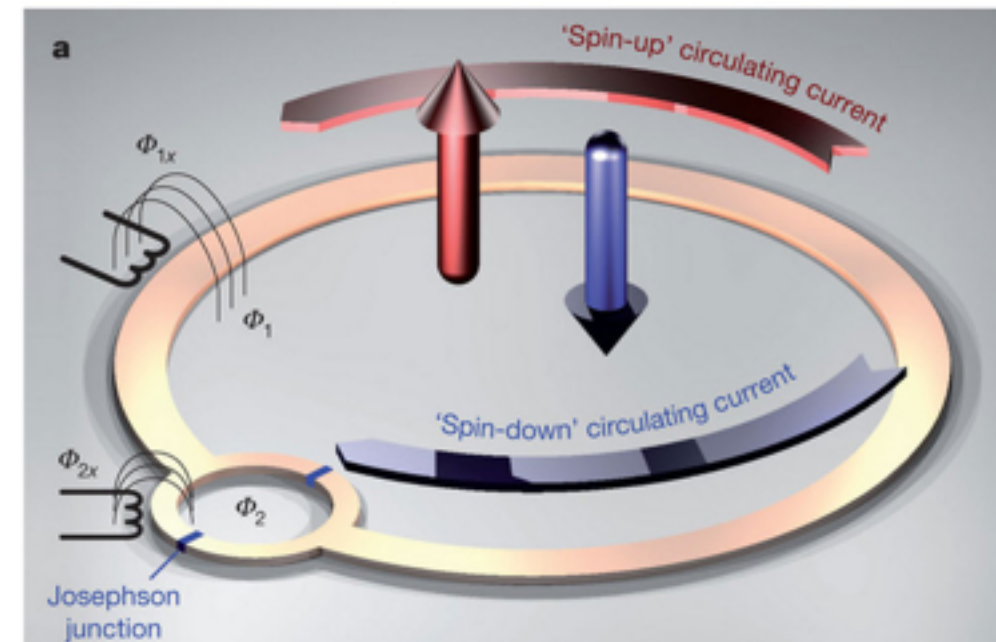
The best qubits: ion traps (universities)

- Use the motional states of well isolated ions to encode a qubit
 - up to 14 qubits
 - coherent for several seconds
 - about 100 gate operations
 - 10 μ s gate times
- Advantages and disadvantages
 - Well isolated from environment
 - Relatively slow
 - Hard to scale beyond $O(20)$ qubits



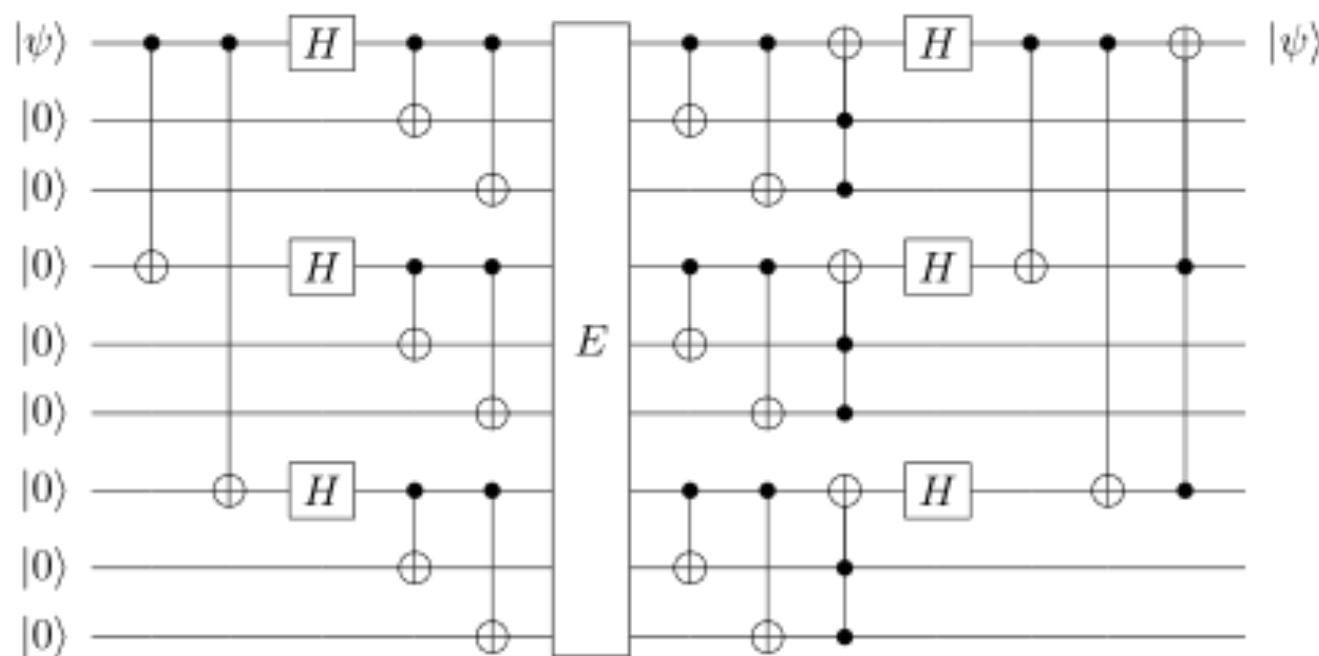
Superconducting qubits (IBM and universities)

- Use superconducting current loops to encode a qubit
- State of the art
 - 5 qubits
 - 100 μ s coherence times
 - 10 ns gate times
 - about 100 gate operations
- Advantages and disadvantages
 - scalable
 - fast
 - can be built in semiconductor foundries
 - lots of decoherence in a solid state device



Quantum error correction

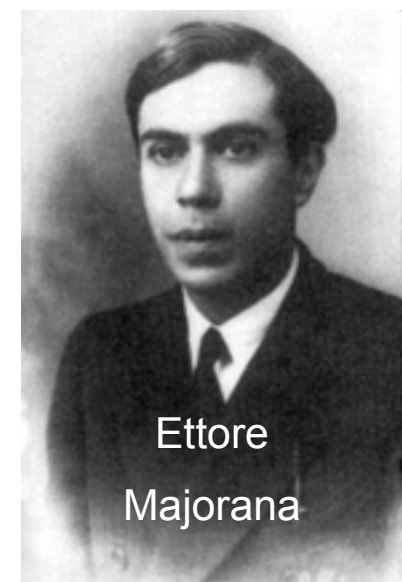
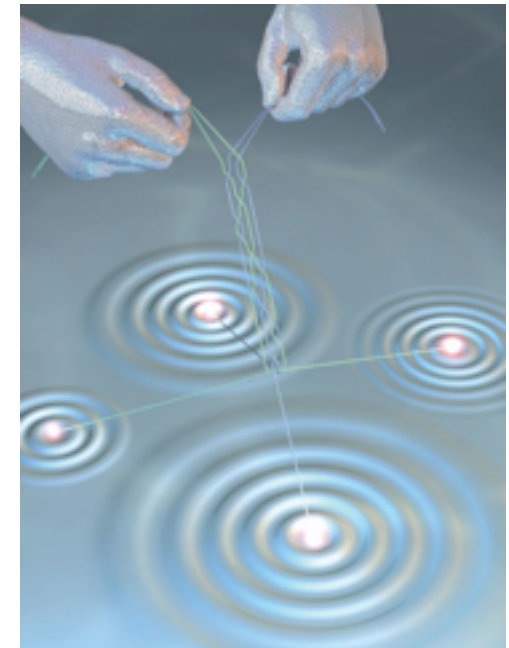
- **Example: Shor's code** protects against sign or bit flip errors
- Needs nine qubits and substantially increases number of gates



- Codes need to be recursively iterated to protect against multiple errors
- Quantum error correction overhead is at least 1000x more qubits and gates for any reasonable computation

Topological qubits (Microsoft and universities)

- Encode the qubit in a topological property of a quantum state
 - Intrinsic protection due to topology
 - No local noise can decohere the qubit
 - No expensive error correction needed
 - Operations done by “braiding” isolated particles
- Most promising:
Majorana particles
 - may exist at ends of superconducting nanowires
 - hope to confirm their existence within 5 years



(Optimistic) road map for quantum computers

- First quantum devices exist, but computational power is limited

- quantum random numbers
- quantum encryption
- quantum simulators
- quantum optimizers



- General-purpose quantum computers will always remain special purpose accelerators for selected tasks
- My optimistic road map
 - Today: 14 qubits allowing up to 100 operations
 - 2020: a long-time stable quantum bit with quantum error correction
 - 2025: matching classical CPUs on selected applications
 - 2030: outperforming any hypothetical classical computer on selected applications
- Huge potential for classical spin-off technology (e.g. superconducting)

Quantum computing applications

- Identifying killer-apps for quantum computing is challenging
 - the problem has to be hard enough that it cannot be solved on a PC
 - the problem has to be amenable to quantum acceleration
 - the crossover scale has to be short enough to make it useful
 - CMOS technology is a tough competitor
 - we need to consider special purpose classical devices as competitors
- Potential applications
 - factoring and code breaking (limited use)
 - quantum chemistry and material simulations (challenging but enormous potential)
 - solving linear systems (can't we solve them well enough already?)
 - machine learning???
- It is time to move quantum computing research from theoretical computer science to high performance computing